

For Wednesday

- Finish chapter 11
- Homework:
 - Chapter 11, exercise 4

Program 3

- Any questions?

Planning in Situation Calculus

$PlanResult(p,s)$ is the situation resulting from executing p in s

$$PlanResult([],s) = s$$

$$PlanResult([a/p],s) = PlanResult(p,Result(a,s))$$

Initial state $At(Home,S_0) \wedge \neg Have(Milk,S_0) \wedge \dots$

Actions as Successor State axioms

$$Have(Milk,Result(a,s)) \Leftrightarrow [(a=Buy(Milk) \wedge At(Supermarket,s)) \vee Have(Milk,s) \wedge a \neq \dots]$$

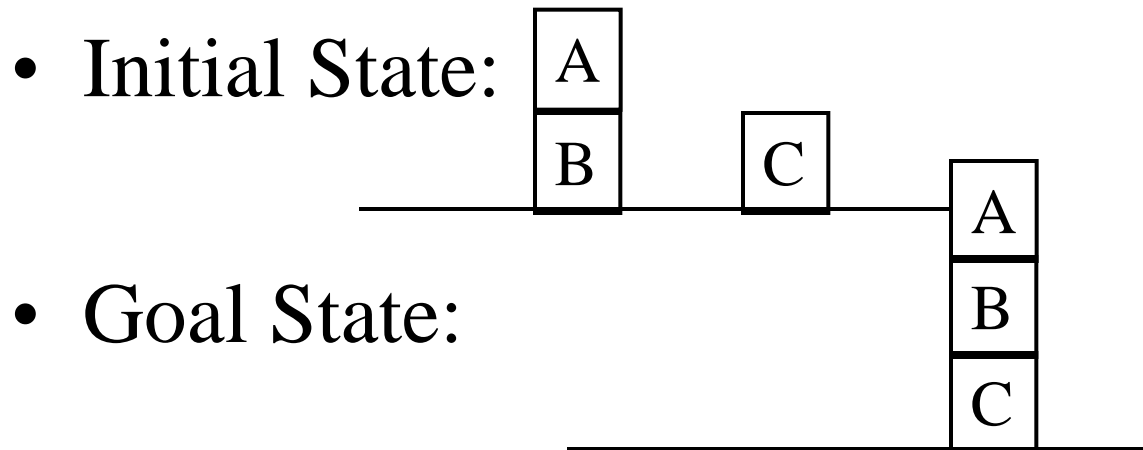
Query $s=PlanResult(p,S_0) \wedge At(Home,s) \wedge Have(Milk,s) \wedge \dots$

Solution $p = Go(Supermarket),Buy(Milk),Buy(Bananas),Go(HWS),\dots]$

- Principal difficulty: unconstrained branching, hard to apply heuristics

The Blocks World

- We have three blocks A, B, and C
- We can know things like whether a block is **clear** (nothing on top of it) and whether one block is **on** another (or on the table)



Situation Calculus in Prolog

```
holds(on(A,B),result(puton(A,B),S)) :-  
    holds(clear(A),S), holds(clear(B),S),  
    neq(A,B).
```

```
holds(clear(C),result(puton(A,B),S)) :-  
    holds(clear(A),S), holds(clear(B),S),  
    holds(on(A,C),S),  
    neq(A,B).
```

```
holds(on(X,Y),result(puton(A,B),S)) :-  
    holds(on(X,Y),S),  
    neq(X,A), neq(Y,A), neq(A,B).
```

```
holds(clear(X),result(puton(A,B),S)) :-  
    holds(clear(X),S), neq(X,B).
```

```
holds(clear(table),S).
```

neq(a,table).

neq(table,a).

neq(b,table).

neq(table,b).

neq(c,table).

neq(table,c).

neq(a,b).

neq(b,a).

neq(a,c).

neq(c,a).

neq(b,c).

neq(c,b).

Situation Calculus Planner

```
plan([],_,_).
```

```
plan([G1|Gs], S0, S) :-
```

```
    holds(G1,S),
```

```
    plan(Gs, S0, S),
```

```
    reachable(S,S0).
```

```
reachable(S,S).
```

```
reachable(result(_,S1),S) :-
```

```
    reachable(S1,S).
```

- However, what will happen if we try to make plans using normal Prolog depth-first search?

Stack of 3 Blocks

holds(on(a,b), s0).

holds(on(b,table), s0).

holds(on(c,table),s0).

holds(clear(a), s0).

holds(clear(c), s0).

| ?- cpu_time(db_prove(6,plan([on(a,b),on(b,c)],s0,S)), T).

S = result(puton(a,b),result(puton(b,c),result(puton(a,table),s0)))

T = 1.3433E+01

Invert stack

holds(on(a,table), s0).

holds(on(b,a), s0).

holds(on(c,b),s0).

holds(clear(c), s0).

?- cpu_time(db_prove(6,plan([on(b,c),on(a,b)],s0,S)),T).

S = result(puton(a,b),result(puton(b,c),result(puton(c,table),s0))),

T = 7.034E+00

Simple Four Block Stack

holds(on(a,table), s0).

holds(on(b,table), s0).

holds(on(c,table),s0).

holds(on(d,table),s0).

holds(clear(c), s0).

holds(clear(b), s0).

holds(clear(a), s0).

holds(clear(d), s0).

| ?- cpu_time(db_prove(7,plan([on(b,c),on(a,b),on(c,d)],s0,S)),T).

S = result(puton(a,b),result(puton(b,c),result(puton(c,d),s0))),

T = 2.765935E+04

7.5 hours!

STRIPS

- Developed at SRI (formerly Stanford Research Institute) in early 1970's.
- Just using theorem proving with situation calculus was found to be too inefficient.
- Introduced STRIPS action representation.
- Combines ideas from problem solving and theorem proving.
- Basic backward chaining in state space but solves subgoals independently and then tries to reachieve any clobbered subgoals at the end.

STRIPS Representation

- Attempt to address the frame problem by defining actions by a precondition, and add list, and a delete list. (Fikes & Nilsson, 1971).
 - Precondition: logical formula that must be true in order to execute the action.
 - Add list: List of formulae that become true as a result of the action.
 - Delete list: List of formulae that become false as result of the action.

Sample Action

- $\text{Puton}(x,y)$
 - Precondition: $\text{Clear}(x) \wedge \text{Clear}(y) \wedge \text{On}(x,z)$
 - Add List: $\{\text{On}(x,y), \text{Clear}(z)\}$
 - Delete List: $\{\text{Clear}(y), \text{On}(x,z)\}$

STRIPS Assumption

- Every formula that is satisfied before an action is performed and does not belong to the delete list is satisfied in the resulting state.
- Although $\text{Clear}(z)$ implies that $\text{On}(x,z)$ must be false, it must still be listed in the delete list explicitly.
- For action $\text{Kill}(x,y)$ must put $\text{Alive}(y)$, $\text{Breathing}(y)$, $\text{Heart-Beating}(y)$, etc. must all be included in the delete list although these deletions are implied by the fact of adding $\text{Dead}(y)$

Subgoal Independence

- If the goal state is a conjunction of subgoals, search is simplified if goals are assumed independent and solved separately (divide and conquer)
- Consider a goal of A on B and C on D from 4 blocks all on the table

Subgoal Interaction

- Achieving different subgoals may interact, the order in which subgoals are solved in this case is important.
- Consider 3 blocks on the table, goal of A on B and B on C
- If do `puton(A,B)` first, cannot do `puton(B,C)` without undoing (clobbering) subgoal: `on(A,B)`

Sussman Anomaly

- Goal of A on B and B on C
- Starting state of C on A and B on table
- Either way of ordering subgoals causes clobbering

STRIPS Approach

- Use resolution theorem prover to try and prove that goal or subgoal is satisfied in the current state.
- If it is not, use the incomplete proof to find a set of **differences** between the current and goal state (a set of subgoals).
- Pick a **subgoal** to solve and an **operator** that will achieve that subgoal.
- Add the precondition of this operator as a new goal and **recursively** solve it.

STRIPS Algorithm

STRIPS(init-state, goals, ops)

Let current-state be init-state;

For each goal in goals do

 If goal cannot be proven in current state

 Pick an operator instance, op, s.t. goal \in adds(op);

 /* Solve preconditions */

 STRIPS(current-state, preconds(op), ops);

 /* Apply operator */

 current-state := current-state + adds(op) - dels(ops);

 /* Patch any clobbered goals */

Let rgoals be any goals which are not provable in current-state;

STRIPS(current-state, rgoals, ops).

Algorithm Notes

- The “pick operator instance” step involves a nondeterministic choice that is backtracked to if a dead-end is ever encountered.
- Employs chronological backtracking (depth-first search), when it reaches a dead-end, backtrack to last decision point and pursue the next option.

Norvig's Implementation

- Simple propositional (no variables) Lisp implementation of STRIPS.
#S(OP ACTION (MOVE C FROM TABLE TO B)
PRECONDS ((SPACE ON C) (SPACE ON B) (C ON TABLE))
ADD-LIST ((EXECUTING (MOVE C FROM TABLE TO B)) (C ON B))
DEL-LIST ((C ON TABLE) (SPACE ON B)))
- Commits to first sequence of actions that achieves a subgoal (incomplete search).
- Prefers actions with the most preconditions satisfied in the current state.
- Modified to try and re-achieve any clobbered subgoals (only once).

STRIPS Results

; Invert stack (good goal ordering)

> (gps '((a on b)(b on c) (c on table) (space on a) (space on table))
'((b on a) (c on b)))

Goal: (B ON A)

Consider: (MOVE B FROM C TO A)

Goal: (SPACE ON B)

Consider: (MOVE A FROM B TO TABLE)

Goal: (SPACE ON A)

Goal: (SPACE ON TABLE)

Goal: (A ON B)

Action: (MOVE A FROM B TO TABLE)

Goal: (SPACE ON A)

Goal: (B ON C)

Action: (MOVE B FROM C TO A)

Goal: (C ON B)

Consider: (MOVE C FROM TABLE TO B)

Goal: (SPACE ON C)

Goal: (SPACE ON B)

Goal: (C ON TABLE)

Action: (MOVE C FROM TABLE TO B)

((START)

(EXECUTING (MOVE A FROM B TO TABLE))

(EXECUTING (MOVE B FROM C TO A))

(EXECUTING (MOVE C FROM TABLE TO B)))

; Invert stack (bad goal ordering)

> (gps '((a on b)(b on c) (c on table) (space on a) (space on table))
'((c on b)(b on a)))

Goal: (C ON B)

Consider: (MOVE C FROM TABLE TO B)

Goal: (SPACE ON C)

Consider: (MOVE B FROM C TO TABLE)

Goal: (SPACE ON B)

Consider: (MOVE A FROM B TO TABLE)

Goal: (SPACE ON A)

Goal: (SPACE ON TABLE)

Goal: (A ON B)

Action: (MOVE A FROM B TO TABLE)

Goal: (SPACE ON TABLE)

Goal: (B ON C)

Action: (MOVE B FROM C TO TABLE)

Goal: (SPACE ON B)

Goal: (C ON TABLE)

Action: (MOVE C FROM TABLE TO B)

Goal: (B ON A)

Consider: (MOVE B FROM TABLE TO A)

Goal: (SPACE ON B)

Consider: (MOVE C FROM B TO TABLE)

Goal: (SPACE ON C)

Goal: (SPACE ON TABLE)

Goal: (C ON B)

Action: (MOVE C FROM B TO TABLE)

Goal: (SPACE ON A)

Goal: (B ON TABLE)

Action: (MOVE B FROM TABLE TO A)

Must reach clobbered goals: ((C ON B))

Goal: (C ON B)

Consider: (MOVE C FROM TABLE TO B)

Goal: (SPACE ON C)

Goal: (SPACE ON B)

Goal: (C ON TABLE)

Action: (MOVE C FROM TABLE TO B)

((START)

(EXECUTING (MOVE A FROM B TO TABLE))

(EXECUTING (MOVE B FROM C TO TABLE))

(EXECUTING (MOVE C FROM TABLE TO B))

(EXECUTING (MOVE C FROM B TO TABLE))

(EXECUTING (MOVE B FROM TABLE TO A))

(EXECUTING (MOVE C FROM TABLE TO B)))

STRIPS on Sussman Anomaly

> (gps '((c on a)(a on table)(b on table) (space on c) (space on b)
(space on table)) '((a on b)(b on c)))

Goal: (A ON B)

Consider: (MOVE A FROM TABLE TO B)

Goal: (SPACE ON A)

Consider: (MOVE C FROM A TO TABLE)

Goal: (SPACE ON C)

Goal: (SPACE ON TABLE)

Goal: (C ON A)

Action: (MOVE C FROM A TO TABLE)

Goal: (SPACE ON B)

Goal: (A ON TABLE)

Action: (MOVE A FROM TABLE TO B)

Goal: (B ON C)

Consider: (MOVE B FROM TABLE TO C)

Goal: (SPACE ON B)

Consider: (MOVE A FROM B TO TABLE)

Goal: (SPACE ON A)

Goal: (SPACE ON TABLE)

Goal: (A ON B)

Action: (MOVE A FROM B TO TABLE)

Goal: (SPACE ON C)

Goal: (B ON TABLE)

Action: (MOVE B FROM TABLE TO C)

Must reachieve clobbered goals: ((A ON B))

Goal: (A ON B)

Consider: (MOVE A FROM TABLE TO B)

Goal: (SPACE ON A)

Goal: (SPACE ON B)

Goal: (A ON TABLE)

Action: (MOVE A FROM TABLE TO B)

((START) (EXECUTING (MOVE C FROM A TO TABLE))

(EXECUTING (MOVE A FROM TABLE TO B))

(EXECUTING (MOVE A FROM B TO TABLE))

(EXECUTING (MOVE B FROM TABLE TO C))

(EXECUTING (MOVE A FROM TABLE TO B)))

How Long Do 4 Blocks Take?

;; Stack four clear blocks (good goal ordering)

> (time (gps '((a on table)(b on table) (c on table) (d on table)(space on a)
(space on b) (space on c) (space on d)(space on table))

'((c on d)(b on c)(a on b))))

User Run Time = 0.00 seconds

((START)

(EXECUTING (MOVE C FROM TABLE TO D))

(EXECUTING (MOVE B FROM TABLE TO C))

(EXECUTING (MOVE A FROM TABLE TO B)))

:: Stack four clear blocks (bad goal ordering)

> (time (gps '((a on table)(b on table) (c on table) (d on table)(space on a)
(space on b) (space on c) (space on d)(space on table))

'((a on b)(b on c) (c on d))))

User Run Time = 0.06 seconds

((START)

(EXECUTING (MOVE A FROM TABLE TO B))

(EXECUTING (MOVE A FROM B TO TABLE))

(EXECUTING (MOVE B FROM TABLE TO C))

(EXECUTING (MOVE B FROM C TO TABLE))

(EXECUTING (MOVE C FROM TABLE TO D))

(EXECUTING (MOVE A FROM TABLE TO B))

(EXECUTING (MOVE A FROM B TO TABLE))

(EXECUTING (MOVE B FROM TABLE TO C))

(EXECUTING (MOVE A FROM TABLE TO B)))

State-Space Planners

- **State-space** (situation space) planning algorithms search through the space of possible states of the world searching for a path that solves the problem.
- They can be based on **progression**: a forward search from the initial state looking for the goal state.
- Or they can be based on **regression**: a backward search from the goals towards the initial state
- STRIPS is an incomplete regression-based algorithm.

Plan-Space Planners

- **Plan-space** planners search through the space of partial plans, which are sets of actions that may not be totally ordered.
- Partial-order planners are plan-based and only introduce ordering constraints as necessary (**least commitment**) in order to avoid unnecessarily searching through the space of possible orderings

Partial Order Plan

- Plan which does not specify unnecessary ordering.
- Consider the problem of putting on your socks and shoes.

Plans

- A **plan** is a three tuple $\langle A, O, L \rangle$
 - A: A set of **actions** in the plan, $\{A_1, A_2, \dots, A_n\}$
 - O: A set of **ordering constraints** on actions $\{A_i \langle A_j, A_k \langle A_1, \dots, A_m \langle A_n\}$. These must be consistent, i.e. there must be at least one total ordering of actions in A that satisfy all the constraints.
 - L: a set of **causal links** showing how actions support each other

Causal Links and Threats

- A **causal link**, $A_p \rightarrow^Q A_c$, indicates that action A_p has an effect Q that achieves precondition Q for action A_c .
- A threat, is an action A_t that can render a causal link $A_p \rightarrow^Q A_c$ ineffective because:
 - $O \cup \{A_p < A_t < A_c\}$ is consistent
 - A_t has $\neg Q$ as an effect

Threat Removal

- Threats must be removed to prevent a plan from failing
- **Demotion** adds the constraint $A_t < A_p$ to prevent clobbering, i.e. push the clobberer before the producer
- **Promotion** adds the constraint $A_c < A_t$ to prevent clobbering, i.e. push the clobberer after the consumer

Initial (Null) Plan

- Initial plan has
 - $A = \{ A_0, A_\infty \}$
 - $O = \{ A_0 < A_\infty \}$
 - $L = \{ \}$
- A_0 (Start) has no preconditions but all facts in the initial state as effects.
- A_∞ (Finish) has the goal conditions as preconditions and no effects.

Example

Op(Action: Go(there); Precond: At(here);

Effects: At(there), \neg At(here))

Op(Action: Buy(x), Precond: At(store), Sells(store,x);

Effects: Have(x))

- A_0 :
 - At(Home) Sells(SM,Banana) Sells(SM,Milk)
Sells(HWS,Drill)
- A_∞
 - Have(Drill) Have(Milk) Have(Banana)
At(Home)

POP Algorithm

- Stated as a **nondeterministic** algorithm where choices must be made. Various search methods can be used to explore the space of possible choices.
- Maintains an **agenda** of goals that need to be **supported** by links, where an agenda element is a pair $\langle Q, A_i \rangle$ where Q is a precondition of A_i that needs supporting.
- Initialize plan to null plan and agenda to conjunction of goals (preconditions of Finish).
- Done when all preconditions of every action in plan are supported by causal links which are not threatened.

POP($\langle A, O, L \rangle$, agenda)

- 1) **Termination**: If agenda is empty, return $\langle A, O, L \rangle$.
Use topological sort to determine a totally ordered plan.
- 2) **Goal Selection**: Let $\langle Q, A_{\text{need}} \rangle$ be a pair on the agenda
- 3) **Action Selection**: Let A_{add} be a nondeterministically chosen action that adds Q . It can be an existing action in A or a new action. If there is no such action return failure.

$$L' = L \cup \{A_{\text{add}} \rightarrow Q, A_{\text{need}}\}$$

$$O' = O \cup \{A_{\text{add}} < A_{\text{need}}\}$$

if A_{add} is new then

$$A' = A \cup \{A_{\text{add}}\} \text{ and } O' = O' \hat{\cup} \{A_0 < A_{\text{add}} < A_{\infty}\}$$

else $A' = A$

4) **Update goal set:**

Let $\text{agenda}' = \text{agenda} - \{ \langle Q, A_{\text{need}} \rangle \}$

If A_{add} is new then for each conjunct Q_i of its precondition,
add $\langle Q_i, A_{\text{add}} \rangle$ to agenda'

5) **Causal link protection:** For every action A_t that threatens a causal link $A_p \rightarrow^Q A_c$ add an ordering constraint by choosing nondeterministically either

(a) **Demotion:** Add $A_t < A_p$ to O'

(b) **Promotion:** Add $A_c < A_t$ to O'

If neither constraint is consistent then return failure.

6) **Recurse:** $\text{POP}(\langle A', O', L' \rangle, \text{agenda}')$

Example

Op(Action: Go(there); Precond: At(there);

Effects: At(there), \neg At(there))

Op(Action: Buy(x), Precond: At(store), Sells(store,x);

Effects: Have(x))

- A_0 :
 - At(Home) Sells(SM,Banana) Sells(SM,Milk)
Sells(HWS,Drill)
- A_∞
 - Have(Drill) Have(Milk) Have(Banana)
At(Home)

Example Steps

- Add three buy actions to achieve the goals
- Use initial state to achieve the Sells preconditions
- Then add Go actions to achieve new preconditions

Handling Threat

- Cannot resolve threat to $At(\text{Home})$ preconditions of both $Go(\text{HWS})$ and $Go(\text{SM})$.
- Must backtrack to supporting $At(x)$ precondition of $Go(\text{SM})$ from initial state $At(\text{Home})$ and support it instead from the $At(\text{HWS})$ effect of $Go(\text{HWS})$.
- Since $Go(\text{SM})$ still threatens $At(\text{HWS})$ of $Buy(\text{Drill})$ must promote $Go(\text{SM})$ to come after $Buy(\text{Drill})$. Demotion is not possible due to causal link supporting $At(\text{HWS})$ precondition of $Go(\text{SM})$

Example Continued

- Add Go(Home) action to achieve At(Home)
- Use At(SM) to achieve its precondition
- Order it after Buy(Milk) and Buy(Banana) to resolve threats to At(SM)